

TRANSMISSION SLOT ALLOCATION METHOD AND MAP FOR VIRTUAL
TUNNELS IN A TRANSMISSION LINE

CROSS-REFERENCE TO RELATED APPLICATIONS

This application is a continuation of U.S. Application Serial No. 09/325,497, filed June 3, 1999 and entitled "TRANSMISSION SLOT ALLOCATION METHOD AND MAP FOR VIRTUAL TUNNELS IN A TRANSMISSION LINE".

5 This application is related to copending U.S. Application Serial No. 09/997,626, filed November 28, 2001 and entitled "METHOD AND SYSTEM FOR TRANSMITTING TRAFFIC IN A VIRTUAL TUNNEL OF A TRANSMISSION LINE", which is a continuation of U.S. Application Serial No. 09/325,687 filed June 3, 1999 and entitled "METHOD AND SYSTEM FOR TRANSMITTING TRAFFIC
10 IN A VIRTUAL TUNNEL OF A TRANSMISSION LINE".

TECHNICAL FIELD OF THE INVENTION

This invention relates generally to the field of telecommunication systems, and more particularly to a transmission slot allocation method and map for virtual tunnels
15 in a transmission line.

BACKGROUND OF THE INVENTION

Telecommunications networks have traditionally been circuit-switch networks that have transmission paths dedicated to specific users for the duration of a call and that employ continuous, fixed-bandwidth transmission. Due to growth in data traffic
5 created by the Internet and its related technologies, however, telecommunications networks are being moved to a packet-switching transmission model. Packet-switch networks provide a large range of digital services, from data to video to basic voice telephony. Packet-switch networks can allow dynamic bandwidth and may be connectionless with no dedicated path or connection-oriented with virtual circuits and
10 dedicated bandwidth along a predetermined path.

Asynchronous transfer mode (ATM) is a connection-oriented packet-switching technology in which information is organized into small, fixed length cells. ATM carries data asynchronously, automatically assigning data cells to available time slots on demand to provide maximum throughput. Compared with
15 other network technologies, ATM provides large increases in maximum supported bandwidth, designed-in isosynchronous traffic support, support for multiple types of traffic such as data, video, and voice transmissions on shared communication lines, and virtual networking capabilities, which increase bandwidth utilization and ease network administration.

20 ATM cells are routed through a telecommunications network at high speeds using a switching label included in the cell. The switching label has two sections that define a virtual path (VP) and a virtual channel (VC) in the network through which the cell is routed. The use of virtual paths and virtual channels allows physical bandwidth in the network to be subdivided and separately commercialized.

25 Because of the low latency and predictability throughput ATM offers, it is capable of providing quality of service (QoS) features. QoS is defined in terms of the attributes of end-to-end ATM connections and is important in an integrated service network, particularly for delay-sensitive applications such as audio and video transmissions, as well as voice-over IP. Other applications in which QoS may be
30 important include traditional data communications, imaging, full-motion video, and multimedia, as well as voice.

Performance criteria for describing QoS for a particular connection include cell loss rate (CLR), cell transfer delay (CTD), and cell delay variation (CDV). ATM traffic is classified as either constant bit rate (CBR) traffic, real-time or non real-time variable bit rate (VBR) traffic, available bit rate (ABR) traffic, and unspecified bit rate (UBR) traffic, depending on the QoS parameters applied to the traffic. CBR and VBR traffic utilize dedicated bandwidth and are intended for real time applications. ABR traffic is intended for non-real time applications which can control, on demand, their transmission rate in a certain range. Like ABR, UBR traffic is intended for non-real time applications which do not have any constraints on the cell delay and cell delay variations.

For CBR, VBR, and other traffic having dedicated bandwidth, transmission slots can be spaced throughout a frame in the transmission line to minimize cell delay variation. However, ABR, UBR, and other types of dynamic bandwidth traffic are evaluated based on transmission line constraints and transmitted in time slots as they become available. As a result, such traffic typically has a high cell delay variation.

SUMMARY OF THE INVENTION

The present invention provides a transmission slot allocation method and map that substantially eliminate or reduce the disadvantages and problems associated with previous systems and methods. In particular, the map stores predefined sets of interleaved port transmission slots that can be assigned to dedicated or dynamic bandwidth traffic.

In accordance with one embodiment of the present invention, a map of transmission slots for a port of a network element includes a plurality of hierarchical sets of port transmission slots. The hierarchical sets include a plurality of parent sets. Each parent set has its port transmission slots divided between a plurality of child sets. The child sets include interleaved port transmission slots.

More specifically, in accordance with a particular embodiment of the present invention, the hierarchical sets each include substantially evenly spaced port transmission slots. In this and other embodiments, each parent set is divided into a same number of child sets. The hierarchical sets may also include a plurality of base sets each having a number of port transmission slots corresponding to a base transmission rate for the port.

Technical advantages of the present invention include providing an improved method and map for allocating bandwidth of a transmission line to a virtual interface. In particular, port transmission slots are grouped into a plurality of hierarchical sets having interleaved port transmission slots. The hierarchical sets can be dynamically allocated to a virtual interface to provide transmission slots for any supported bandwidth. In addition, the transmission slots in each set are substantially evenly spaced throughout the transmission frame of the port. This provides minimal delay variation for both dedicated and dynamic bandwidth traffic, and for multiple virtual interfaces operating at different rates. The hierarchical sets are also sized to be equal to or only slightly larger than a supported transmission rate. As a result, allocatable bandwidth is maximized.

Other technical advantages of the present invention will be readily apparent to one skilled in the art from the following figures, description, and claims.

BRIEF DESCRIPTION OF THE DRAWINGS

For a more complete understanding of the present invention and its advantages, reference is now made to the following description taken in conjunction with the accompanying drawings, wherein like reference numerals represent like parts, in which:

FIGURE 1 is a block diagram illustrating a telecommunications system in accordance with one embodiment of the present invention;

FIGURE 2 is a block diagram illustrating an add/drop multiplexer element for the telecommunications system of FIGURE 1 in accordance with one embodiment of the present invention;

FIGURE 3 is a block diagram illustrating a line card for the add/drop multiplexer of FIGURE 2 in accordance with one embodiment of the present invention;

FIGURE 4 is a schematic diagram illustrating a map of port transmission slots for the line card of FIGURE 3 in accordance with one embodiment of the present invention;

FIGURE 5 is a flow diagram illustrating a method for generating the map of FIGURE 4 in accordance with one embodiment of the present invention;

FIGURE 6 is a flow diagram illustrating a method for selecting evenly spaced transmission slots for generation of the map of FIGURE 4 in accordance with one embodiment of the present invention; and

FIGURE 7 is a flow diagram illustrating the operation of the add/drop multiplexer of FIGURE 2 in accordance with one embodiment of the present invention.

DETAILED DESCRIPTION OF THE INVENTION

FIGURE 1 illustrates a telecommunications system 10 in accordance with one embodiment of the present invention. In this embodiment, the telecommunications system 10 comprises a public portion 12 and a private portion 14 that each transmit voice, data, other suitable types of information, and/or a combination of different types of information. The public portion 12 includes a public network 20 such as the Internet. The private portion 14 includes remote private networks 22 and 24. The private networks 22 and 24 are each an Intranet or other suitable network such as a private local area network (LAN) or a private wide area network (WAN). The telecommunications system 10 may be entirely implemented in either the public network 20 or in one of the private networks 22 or 24, or may be otherwise suitably distributed between disparate networks.

Referring to FIGURE 1, the telecommunications system 10 includes a plurality of nodes 26 interconnected by transmission lines 28. The nodes 26 each comprise a network element capable of communicating traffic in the telecommunications system 10. The network element may be a switch, router, add/drop multiplexer or other suitable device capable of directing traffic in the telecommunications system 10.

The transmission lines 28 provide a physical interface between the nodes 26. Each physical interface is defined by the bandwidth of the connecting transmission line 28 which may be a DS-3 line, an OC-3 line, an OC-12 line, or any other suitable bandwidth. The transmission lines 28 each comprise optical fiber, cable, twisted wire, or other suitable wireline or wireless link capable of transporting traffic between two nodes 26.

For the embodiment of FIGURE 1, the telecommunications system 10 includes a private boundary node 30 ("Node A") at a boundary of the private network 22, a public boundary node 32 ("Node B") at a first boundary of the public network 20, an intermediate node 34 within the public network 20, a public boundary node 36 ("Node C") at a second boundary of the public network 20, and a private boundary node 38 ("Node D") at a boundary of the private network 24. The nodes 30, 32, 34, 36, and 38 are each asynchronous transport mode (ATM) elements or multi-protocol elements having functionality for processing ATM traffic in which information is

organized into small, fixed length cells. The ATM cells are routed by the nodes 30, 32, 34, 36, and 38 using switching labels included in the cell. The switching label has two sections that define a virtual path (VP) and a virtual channel (VC) in the telecommunications system 10 through which the cell is routed. Use of virtual paths
5 (VPs) and virtual channels (VCs) allows physical bandwidth of the telecommunications system 10 to be subdivided and separately commercialized.

A virtual channel (VC) 40 is formed in the transmission line 28 between Nodes A and B. The virtual channel (VC) 40 forms a public universal network interface (UNI) 42 between Node A in the private network 22 and Node B in the
10 public network 20. Similarly, a virtual channel (VC) 44 is formed in the transmission line 28 between Nodes C and D. The virtual channel (VC) 44 forms a public universal network interface (UNI) 46 between Node C in the public network 20 and Node D in the private network 24. Virtual channels (VCs) 40 and 44 together with virtual channels (VCs) 48 in the public network 20 between Nodes B and C form a
15 virtual tunnel 50 between Nodes A and D. The virtual tunnel 50 is a virtual channel (VC), virtual path (VP), or any other suitable portion and/or construct of one or more transmission lines 28. The virtual tunnel 50 together with a virtual interface 52 at the entrance to the virtual tunnel 50 and a termination interface 54 at the end of the virtual tunnel 50 forms a virtual network network interface (V-NNI) between Nodes A and
20 D. Nodes B and C need not know of the virtual network network interface (V-NNI). The virtual interface 52 aggregates and transmits traffic designating the virtual tunnel 50. Traffic designates the virtual tunnel 50 when the traffic includes information indicating that it is to be transmitted in the virtual tunnel 50. The termination interface 54 separates traffic received from the virtual tunnel 50 for individualized
25 routing. Accordingly, virtual channel connections (VCCs) 56 can be transmitted through the telecommunications system 10 in the virtual tunnel 50 and need not compete with other designated traffic outside of the virtual tunnel 50 and inside the system 10.

FIGURE 2 illustrates details of Node A in accordance with one embodiment
30 of the present invention. In this embodiment, Node A is a multi-protocol add/drop multiplexer 80. As described in more detail below, the add/drop multiplexer 80 includes a plurality of transmission line interfaces that each independently queue and

schedule traffic for a transmission line or a virtual tunnel in a transmission line. Remaining Nodes B, C, D, and E may be similarly configured.

Referring to FIGURE 2, the add/drop multiplexer 80 comprises a service complex 82, a transport complex 84, switch complex 86, a management complex 88,
5 and a synchronization complex 90. The disparate complexes 82, 84, 86, 88, and 90 are interconnected by one or more networks or backplanes in the add/drop multiplexer 80.

The service complex 82 includes a plurality of tributary cards 92. The tributary cards 92 each receive customer input/output (I/O) and perform conversions
10 necessary for processing by the switch complex 86.

The transport complex 84 includes a plurality of line cards 94 for transmitting data on the transmission lines 95. In a particular embodiment, the line cards 94 are OC-3 or OC-12 line cards that are bi-directional and can handle optical to electrical (O/E) and electrical to optical (E/O) conversions. In this embodiment, the line cards
15 94 also handle the line and selection overhead extraction and insertion.

For the embodiment of FIGURE 2, the transport complex 84 also includes an interface controller 96 and an interface array 98 comprising a plurality of transmission line interfaces 100. The interface controller 96 is implemented by software stored in a computer-readable medium for execution by a processor of the add/drop multiplexer
20 80. The computer-readable medium is a random access memory (RAM), a read only memory (ROM), or other suitable medium capable of storing instructions for execution by a processor. The interface array 98 is implemented by an application specific integrated chip (ASIC), software, or a combination of hardware and software. Both the interface controller 96 and interface array 98 may, in another embodiment,
25 be separately implemented in each of the line cards 94, or for each port of the line cards 94.

As described in more detail below, the interface controller 96 is operable to generate a port memory map for each port of the transport complex 84 and to couple transmission line interfaces 100 from the interface array 98 to the ports of the line
30 cards 94. A transmission line interface 100 is coupled to a port when it is connected to the port for use as a physical interface in connection with a transmission line 95 or as a virtual interface in connection with a virtual tunnel in the transmission line 95.

An interface is physical when it is for a physical rate of a port or virtual when for a virtual rate of the port. The transmission line interface 100 is coupled to the port by a dedicated or other suitable bus or by dedicated or other suitable lines and/or switches. The interface controller 96 is also operable to selectively allocate port transmission
5 slots from the port memory map to the transmission line interfaces 100 coupled to the port. Accordingly, the interface controller 96 can dynamically implement and size virtual interfaces in each of the outlet ports of the add/drop multiplexer 80.

The switch complex 86 comprises a plurality of switch fabric cards including an ATM switch fabric card. The ATM switch fabric card receives ATM cells on an
10 input port and switches them to an output port. In switching the ATM cells, the switch complex 86 first translates necessary virtual path (VP) and virtual channel (VC) addresses.

The management complex 88 monitors and controls the status of the service, transport, switch, and synchronization complexes 82, 84, 86, and 90. The
15 management complex 88 also maintains alarm, protection switching, and provisioning databases for the add/drop multiplexer 80. The synchronization complex 90 synchronizes the service, transport, and switch complexes 82, 84, and 86 by providing a stable traceable reference clock.

FIGURE 3 illustrates an ATM switch fabric 102 for use in the add/drop
20 multiplexer 80 in accordance with one embodiment of the present invention. In this embodiment, the ATM switch fabric 102 comprises a plurality of ATM output ports 104 for connection to disparate transmission lines 106. The outlet ports 104 each receive ATM traffic from the switch complex 86 of the add/drop multiplexer 80 and transmit the traffic in a corresponding transmission line 106. The ATM traffic
25 comprises dedicated traffic such as constant bit rate (CBR) traffic and real-time and non real-time variable bit rate (VBR) traffic, as well as dynamic traffic such as available bit rate (ABR) traffic and unspecified bit rate (UBR) traffic. Dynamic traffic is traffic that is subject to change, has bandwidth that can be shared between connections, and/or has an unspecified rate within a range. The constant bit rate
30 (CBR), real-time variable bit rate (VBR), and other dedicated traffic utilize substantially constant and/or dedicated bandwidth and are intended for real time applications such as audio, video, and voice-over IP transmissions. The dynamic

traffic is generally bursty in nature. The available bit rate (ABR) traffic is intended for real time applications which can control, on demand, their transmission rate in a certain range. The unspecified bit rate (UBR) traffic is intended for non-real time applications which do not have any constraints on cell delay and cell delay variation.

5 Quality of service (QoS) levels are defined for each traffic type based on cell loss rate (CLR), cell transfer delay (CTD), and cell delay variation (CDV).

Referring to FIGURE 3, the outlet ports 104 each include one or more interfaces 114 and 118 and a port memory map 112. Each interface 114 or 118 is defined by a transmission line interface 100 of the interface array 98 coupled to the

10 port 104. The transmission line interfaces 100 form virtual interfaces 114 for virtual tunnels 116 in the transmission line 106 and physical interfaces 118 for the physical bandwidth of the transmission line 106. The virtual interfaces 114 each transmit traffic designating a corresponding virtual tunnel 116 that is defined by a portion of the bandwidth of the transmission line 106. The physical interfaces 118 each transmit

15 traffic of the port 104 in the bandwidth of the transmission line 106.

The virtual and physical interfaces 114 and 118 each comprise a queue 120 and a scheduler 122. The queue 120 receives designated virtual channel connections (VCC) of any designated class and temporarily stores the traffic for the scheduler 122. The scheduler 122 shapes traffic from the queue 120 and transmits the traffic in the

20 transmission line 106 or the virtual tunnel 116 of the transmission line 106. In shaping the traffic, the scheduler 122 aggregates virtual channel connection (VCC) traffic in a corresponding transmission line 106 and virtual tunnel 116 and ensures that the sum of aggregate virtual channel connection (VCC) traffic does not exceed the transmission line rate of the corresponding virtual tunnel 116. The virtual and

25 physical interfaces 114 and 118 may also include back pressure transmission rate adaptation buffers to account for the egress rate of the scheduler 122 not exactly matching the physical link rate due to limited schedule table length and the like.

For a virtual interface 114, the scheduler 122 transmits traffic from the queue 120 in a set of port transmission slots allocated to the interface 114 from the port

30 memory map 112. The allocated port transmission slots define the virtual tunnel 116 corresponding to the virtual interface 114. The capacity of the virtual interface is equal to the transmission line rate of the virtual tunnel 116. The scheduler 122

ensures the transmission line rate allocated to the virtual tunnel 116 is not exceeded. For the physical interfaces 118, the scheduler 122 transmits traffic from the queue 120 in the port transmission slots representing physical bandwidth. Accordingly, the physical interface 118 uses the physical bandwidth of the corresponding transmission
5 line 106.

For the embodiment of FIGURE 3, the bandwidth of a first transmission line 106 is entirely subdivided into a plurality of virtual tunnels 116. For example, the first transmission line 106 may be an OC-3c transmission line subdivided into three disparate virtual tunnels 116, comprising transmission rates of 100 Mbps, 35 Mbps,
10 and 20 Mbps, respectively. To support this configuration, the port 104 includes separate virtual interface 114 for each virtual tunnel 116. The bandwidth of a second transmission line 132 is subdivided into a virtual tunnel 116 and a remaining portion 134 that transmits traffic of the port 104 not designating the virtual tunnel 116. For example, the second transmission line 132 may be an OC-3c transmission line
15 comprising a 20 Mbps virtual tunnel 116. To support this configuration, the port 104 includes separate virtual interfaces 114 for the virtual tunnel 116 and the remaining bandwidth of the transmission line 106. The interface of the remaining bandwidth is illustrated as a physical interface to illustrate that concept.

FIGURE 4 illustrates the details of a port memory map 140 for each port 104
20 of the ATM switch fabric 102 in accordance with one embodiment of the present invention. In this embodiment, transmission slots of a port are grouped based on a base transmission rate for the port. The base transmission rate is a foundation rate for the port 104 or virtual tunnels 116 supported by the port 104. The base transmission should provide the desired granularity for assignable bandwidth. The base
25 transmission rate is preferably the lowest transmission rate of the port or of a virtual tunnel supported by the port and may be a standardized rate such as the DS-1 rate or other suitable rate.

Referring to FIGURE 4, port memory map 140 comprises a plurality of hierarchal sets 142 of port transmission slots. The hierarchical sets 142 are arranged
30 in a series of hierarchical levels 145 with a plurality of hierarchical sets 142 in one level 145 together containing the port transmission slots of a hierarchical set 142 in a next higher level 145. The port transmission slots are the transmission slots available

to the port 104 for transmission of traffic in a connected transmission line 106. Accordingly, the number of port transmission slots varies with the bandwidth of the transmission line 106. For example, a port 104 adapted for connection to an OC-3 transmission line 106 comprises 4,096 port transmission slots while a port 104
5 adapted for connection to an STS-1 transmission line 106 comprises 1,365 transmission slots.

The hierarchal sets 142 of port transmission slots include a plurality of parent sets 144. Each parent set 144 has its port transmission slots equally divided between a plurality of child sets 146, which may themselves be parent sets 144 to further child
10 sets 146. The port transmission slots of the parent sets 144 are interleaved between the child sets 146 to maintain substantially equal spacing between the transmission slots in each set 144 and 146. Unless otherwise specified, the term each means every one of at least a subset of the identified items. For a parent set 144 having two child sets 146, for example, each child set 146 preferably has every other port transmission
15 slot of the parent set 144.

For the embodiment of FIGURE 4, the port memory map 140 comprises 4,096 port transmission slots for an OC-3 port. Each parent set 144 is divided into two child sets 146. The base sets 148 each comprise forty-one (41) port transmission slots corresponding to a base DS-1 transmission rate. The number of transmission slots
20 correspond to a transmission rate when the number of transmission slots is substantially a minimum number necessary to support the rate. The substantial minimum number necessary to support the rate is the smallest integer number of slots that results in a transmission rate equal to or slightly larger than the actual rate.

The hierarchical sets 142 are segregated into primary hierarchical sets 150 and
25 secondary hierarchical sets 152. The primary hierarchical sets 150 comprise a maximum number of port transmission slots that can be recursively divided evenly to reach the base set 148. For the base DS-1 sets 148 in the illustrated embodiment, the primary hierarchical sets 150 comprise 2,624 port transmission slots. The secondary hierarchical sets 152 comprise the remaining 1,472 port transmission slots of the
30 OC-3 bandwidth.

The primary hierarchical sets 150 comprise an initial set (P_{64}) at a first primary level 145. The initial set (P_{64}) includes all of the 2,624 port transmission slots of the primary hierarchical sets 150. As described in more detail below, these port transmission slots are substantially equally spaced throughout the port transmission frame to provide minimal cell delay variation (CDV). The initial set (P_{64}) is divided into two child sets (P_{32-1} and P_{32-2}) each having a disparate interleaved half of the port transmission slots of the initial set (P_{64}). Accordingly, the port transmission slots in the child sets (P_{32-1} and P_{32-2}) remain substantially equally spaced. The child sets (P_{32-1} and P_{32-2}) are divided into further child sets ($P_{16-1} \dots P_{16-4}$) which are recursively divided into further child sets ($P_{8-1} \dots P_{8-8}$, $P_{4-1} \dots P_{4-16}$, $P_{2-1} \dots P_{2-32}$, and $P_{1-1} \dots P_{1-64}$) until base transmission sets ($P_{1-1} \dots P_{1-64}$) are reached. The resulting primary sets 150 each have substantially evenly spaced transmission slots and represent a particular bandwidth in the transmission line. Accordingly, a primary set 150 can be allocated to a virtual interface 114 to provide desired bandwidth with minimal cell delay variation (CDV).

The secondary hierarchical sets 152 comprise a plurality of tiers 154 each comprising a maximum number of remaining port transmission slots that can be recursively divided evenly to reach the base set 148. The port transmission slots in each tier 154 are evenly spaced to the extent possible to minimize cell delay variation (CDV). In each tier 154, as previously described in connection with the primary hierarchical sets 150, an initial set (S_{32} , S_2 , or S_1) is recursively divided into two child sets 146 until the base set 148 is reached. Each child set 146 comprises an interleaved half of the port transmission slots of the parent set 144.

Each hierarchical set 142, whether parent or child 144 or 146 or primary or secondary 150 or 152 can be allocated to a scheduler 122 of a transmission line interface 100 coupled to the port 104. To optimize cell delay variation for the interface 100, each is preferably assigned a single hierarchical set 142. However, multiple hierarchical sets 142 can be assigned as needed to optimize bandwidth usage.

FIGURE 5 is a flow diagram illustrating a method for generating the port memory map 140 for a port 104 in accordance with one embodiment of the present

invention. In this embodiment, the port 104 has a number of transmission slots that cannot be recursively divided evenly to yield the base set 148. Accordingly, the port memory map 140 will include both primary and secondary sets 150 and 152. In other embodiments in which the total number of port transmission slots can be recursively and evenly divided to reach the base set 148, all resulting sets may be primary sets 150 with optimized spacing. Also in this embodiment, the port memory map 140 is generated by the interface controller 96 in the add/drop multiplexer 80. The port memory map 140 may be independently generated and provided to the add/drop multiplexer 80 for the port 104.

Referring to FIGURE 5, the method begins at step 160 in which the number of port slots in a base set 148 is obtained. The number of port slots correspond to a base transmission rate for the port 104. As previously described, the base transmission rate is the foundation rate for the port 104 or a virtual tunnel 116 supported by the port 104.

Proceeding to step 162, the number of port transmission slots for the primary hierarchical sets 150 is determined. This number comprises a maximum number of port transmission slots that can be recursively divided evenly to reach the base set 148. For the embodiment of FIGURE 4 in which each parent set 144 is divided into two child sets 146, the number of port transmission slots for the primary hierarchical sets 150 is the largest $2N$ multiple of the number of port transmission slots in a base set 148 that is less than or equal to the total number of port transmission slots.

Next, at step 164, the determined number of port transmission slots for the primary hierarchical sets 150 are selected as an initial set from the entirety of the port transmission slots. The selected port transmission slots are substantially evenly spaced to provide minimal cell delay variation (CDV) for the primary hierarchical sets 150. Further details for selecting substantially evenly spaced port transmission slots are described below in connection with FIGURE 6.

At step 166, the initial set is divided into two child sets 146. Each child set 146 has a disparate interleaved half of the port transmission slots of its parent set 144. Next, at step 168, each child set 146 is recursively divided into two lower-level 145 child sets 146 until the child sets 146 in a lowest level 145 are base sets 148.

Accordingly, each of the primary hierarchical sets 150 correspond to a base rate or is a multiple of the base rate.

Proceeding to step 170, a number of port transmission slots for a next tier 154 of secondary hierarchical sets 152 is determined. The next tier secondary hierarchical sets 152 comprise a maximum number of remaining port transmission slots that can be recursively divided evenly to reach the base set 148. The remaining port transmission slots are transmission slots of the port 104 not previously selected for the primary hierarchical sets 150 or a previous tier 154 of secondary hierarchical sets 152.

At step 172, the determined numbered port transmission slots for the next tier secondary hierarchical sets 152 are selected as an initial set from the remaining port transmission slots. The selected port transmission slots are evenly spaced in the port transmission frame to the extent possible to minimize cell delay variation (CDV) for the next tier secondary hierarchical sets 152. Further details for selecting the port transmission slots are described below in connection with FIGURE 6.

Proceeding to step 174, the initial set is divided into two child sets 146. Each child set 146 has a disparate interleaved half of the port transmission slots of its parent set 144. At step 176, each child set 146 is recursively divided into two lower-level 145 child sets 146 until the child sets in a lowest level 145 are base sets 148. Accordingly, each secondary hierarchical set 152 correspond to a base transmission rate or is a multiple of the base transmission rate.

Proceeding to decisional step 178, the interface controller 98 determines whether the number of unselected port transmission slots are capable of supporting additional base sets 148. If the unselected port transmission slots are capable of supporting additional base sets 148, the remaining port transmission slots can be further grouped into hierarchical sets 142 and the Yes branch of decisional step 178 returns to step 170 where the number of port transmission slots for the next tier of secondary hierarchical sets 152 is determined. After all port transmission slots capable of supporting a base set 148 have been selected, the No branch of decisional step 178 leads to step 180. At step 180, the remaining port transmission slots are selected as a remainder set that can be combined with one or more of the hierarchical sets 142 to increase the bandwidth of a virtual tunnel 116, be used as a limited

non-standard virtual tunnel 116, or be otherwise suitably used as needed to transport data from the port 104.

Next, at step 182, the initial, child, and remainder sets are saved to a computer readable medium as the port memory map 140. In this way, port transmission slots
5 are grouped into a plurality of hierarchical sets 142 having interleaved port transmission slots. The hierarchical sets 142 can be dynamically allocated to virtual or physical interfaces 114 or 118 to provide transmission slots for any supportive bandwidth. The hierarchical sets 142 are also sized to be equal to or only slightly larger than a base transmission rate. As a result, allocatable bandwidth in a
10 transmission line 106 is maximized.

FIGURE 6 is a flow diagram illustrating a method for selecting substantially evenly spaced transmission slots for the primary and secondary sets 150 and 152 of the port memory map 140 in accordance with one embodiment of the present invention. The method may be used to map any lower rate map of transmission slots
15 into a higher rate map of transmission slots.

Referring to FIGURE 6, the method begins at step 200 in which a number of port transmission slots to be selected for a lower rate map is obtained. For the port memory map 140, the lower rate map may be the primary hierarchical sets 150 or a tier of secondary hierarchical sets 152.

Proceeding to step 202, the number of port transmission slots available in a
20 higher rate map for allocation to the lower rate map is obtained. For the primary hierarchical sets 150, the number of port transmission slots available for allocation to the primary hierarchical sets 150 comprise the total number of transmission slots for the port 104. Accordingly, spacing may be optimized for the primary hierarchical set
25 150. For a tier of secondary hierarchical sets 152, the number of port transmission slots available for allocation to the secondary hierarchical sets 152 comprise the remaining transmission slots of the port that have not been previously allocated to the primary hierarchical sets 150 or previous tiers of secondary hierarchical sets 152.

Next, at step 204, an average slot interval is obtained. As used herein,
30 obtained means determined or received. In one embodiment, the average slot interval comprises the ratio of the number of port transmission slots available for selection in

the higher rate map divided by the number of port transmission slots to be selected for the lower rate map.

Proceeding to step 206, a starting position in the higher rate map is obtained. In one embodiment, the starting position is the first available port transmission slot in the higher rate map. At step 208, the port transmission slot at the starting position is
5 selected as the first port transmission slot for the lower rate map.

At step 210, an estimated position of a next port transmission slot is generated by incrementing the position of the previous port transmission slot by the average slot interval. Accordingly, port transmission slots for the lower rate map are each
10 separated by the average interval to provide optimized cell delay variation. At step 212, the next port transmission slot is selected based on the estimated position. In one embodiment, a floor position is determined from a floor function of the estimated position. In this embodiment, the transmission slot at the floor position is selected as the next port transmission slot.

Proceeding to decisional step 214, it is determined whether all of the port transmission slots have been selected for the lower rate map. If all of the port transmission slots have not yet been selected, the No branch of decisional step 214 returns to step 210 where an estimated position of the next port transmission slot is generated. After all of the port transmission slots have been selected, the Yes branch
15 of decisional step 214 leads to the end of the process. In this way, port transmission slots are selected by the interface controller 96 or other device during generation of the port memory map 104 or for other lower rate maps in a higher rate map. The transmission slots are substantially evenly spaced to the extent possible throughout the transmission frame of the port. This provides minimal cell delay variation for
20 both dedicated and dynamic bandwidth traffic, and for multiple virtual interfaces operating at different rates.

FIGURE 7 is a flow diagram illustrating operation of the add/drop multiplexer 80 in accordance with one embodiment of the present invention. In this embodiment, the add/drop multiplexer 80 processes ATM traffic including a plurality of virtual
30 channel connections (VCCs) designating a virtual tunnel 116 in a transmission line 106 connected to an outlet port 104 of the multiplexer 80. The virtual tunnel 116 is

controlled by a virtual interface 114 and defined by a hierarchical set 142 of port transmission slots allocated to the virtual interface 114.

Referring to FIGURE 7, the method begins at step 250 in which the add/drop multiplexer 80 receives disparate virtual channel connections (VCCs) designating the
5 virtual tunnel 116 from one or more sources. At step 252, the add/drop multiplexer 80 translates necessary virtual path (VP) and virtual channel (VC) addresses for the virtual channel connections (VCCs). Based on the translated address, the add/drop multiplexer 80 switches the VCCs to the virtual interface 114 for the virtual tunnel 116.

10 Proceeding to step 256, the virtual channel connections (VCCs) are aggregated in the queue 120 of the virtual interface 114. The virtual channel connections (VCCs) may include dedicated traffic such as constant bit rate (CBR) or real-time variable bit rate (VBR) traffic or dynamic traffic such as available bit rate (ABR) traffic or unspecified bit rate (UBR) traffic. At step 258, the scheduler 122 for the virtual
15 interface 114 schedules the virtual channel connections (VCCs) for transmission within the port transmission slots of the allocated hierarchical set 142. The port transmission slots define the virtual tunnel 116. Accordingly, the virtual channel connections (VCCs) need not compete with other traffic in the transmission line 106.

Proceeding to step 260, scheduled traffic is transmitted in the virtual tunnel
20 116 of the transmission line 106. The traffic may travel in the virtual tunnel 116 to a next node or through a plurality of intermediate nodes to a remote termination node. At step 262, the virtual channel connections (VCCs) are separated and individually routed. In this way, dynamic traffic is transported within a virtual tunnel, and can even be combined with dedicated or other dynamic traffic sources. As a result, virtual
25 tunnels can be used for all types of traffic and quality of service (QoS) levels can be provided for dynamic traffic. This allows a telecommunications provider to better package resources and to better support its customers whose dynamic traffic no longer need compete with that of other businesses.

Although the present invention has been described with several embodiments,
30 various changes and modifications may be suggested to one skilled in the art. It is intended that the present invention encompass such changes and modifications as fall within the scope of the appended claims.